MINE TO PLANT PRODUCTION SCHEDULING

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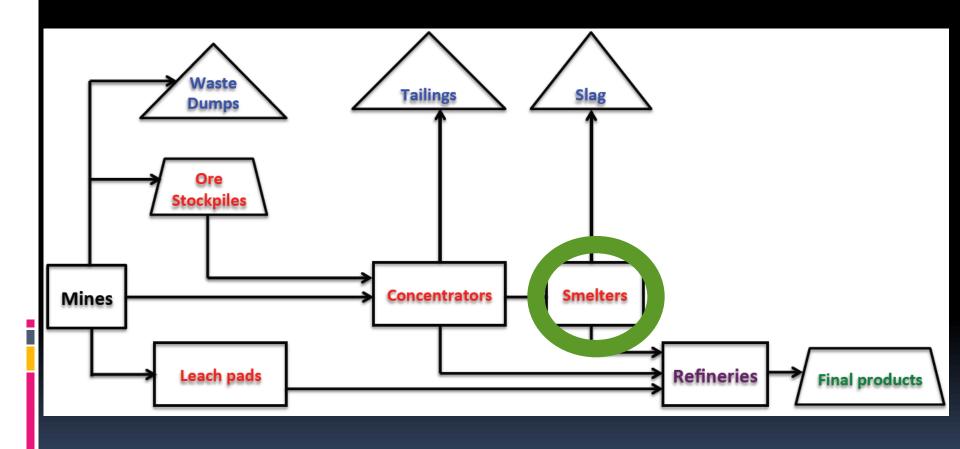
- Classical themes in <u>industrial engineering</u>
 - Production planning
 - Downstream logistics
- Mineral/metallurgical extraction
 - Particular structures
 - More than just plugging numbers into models
- Judgment and expertise to link the two
 - What to solve
 - How to solve

Part 1:

 Strategic metallurgical production planning under geostatistical uncertainty

Part 2:

Short-term smelter production scheduling
 (Automatic scheduling of <u>Altonorte operations</u> using greedy algorithms)



Part 1:

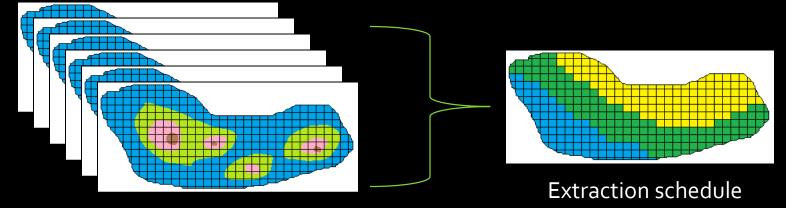
STRATEGIC METALLURGICAL PRODUCTION PLANNING UNDER GEOLOGICAL UNCERTAINTY

Overview

- Stochastic mine planning
- Centralized vs localized decision-making
- Bi-level approach to <u>downstream optimization</u>

Stochastic mine planning

- Capture geostatistical uncertainty
 - Repeated conditional simulations



Orebody simulations

Constrained optimization

 Maximize NPV
 (incomplete picture)
 Respect constraints (capacity, blending, etc.)

 Extend scope of optimization??

 (centralization of decision-making)

Centralized vs localized DM

- Typical concern in industrial engineering
- Balance two perspectives
 - 1. Globalized solutions bring gains

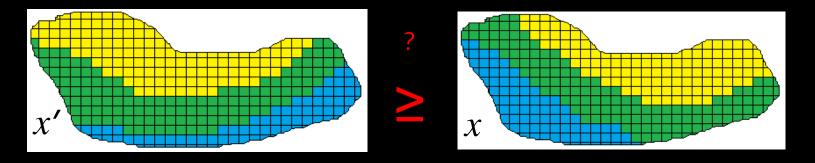
$$\max_{x} \left[f(x, y_{o}) \right] \leq \max_{x, y} \left[f(x, y) \right]$$

- 2. Local decision-makers to react to changing circumstances ("local manageability")
- Low hanging fruit: Reengineering balance between centralization and localization
- Unmanaged variability

 overengineering of downstream operations

Metaheuristics

- Current application of metaheuristics
 - Compare two potential extraction plans



- Compare f(x') and f(x)
- Update data structures accordingly
- Continue searching for improvements...
- <u>Downstream operations</u> may be integrated within objective function *f*.

Bi-level optimization $f(x) = \max_{y} \left[f(x,y) \right]$

An <u>optimization</u> within an <u>optimization</u>

Current approach

f(*x*) is NPV of *x f*(*x*') is NPV of *x*'

Generic allocation of dowstream resources

Overengineered (suboptimal)

Bi-level approach

- f(x) is NPV of x, given an <u>optimal</u> allocation of downstream resources for y(x)
- f(x') is NPV of x', given <u>optimal</u> allocation of downstream resources for y(x')

Bi-level optimization

- Outer optimization
 - Vast solution space
 - (discrete block orderings)
 - Subject to geological uncertainty

Inner optimization

- Continuous solution space
- Dominated by mass flows
- Subject to geological uncertainty

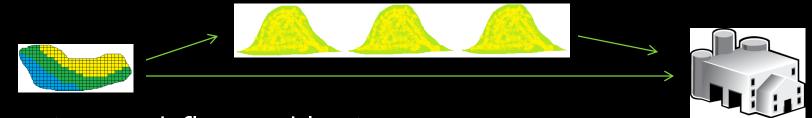
Linear program (with embedded flow network)

Metaheuristics

Two proposals for bi-level formulations

Role of Linear Program

- Automate decisions for stockpiling v/s processing
 - When to <u>send material</u> to stockpiles?
 - When to <u>draw material</u> from stockpiles?
 - (Classical industrial engineering theme)



- (Network flow problem)
- Allocation of resources
 - Divide <u>plant time</u> between several <u>modes of operation</u>
 - Divide transport routes between several product streams
- Evaluate proposed mine extraction plan (f)

Second Proposal

<u>Preprocessing</u>

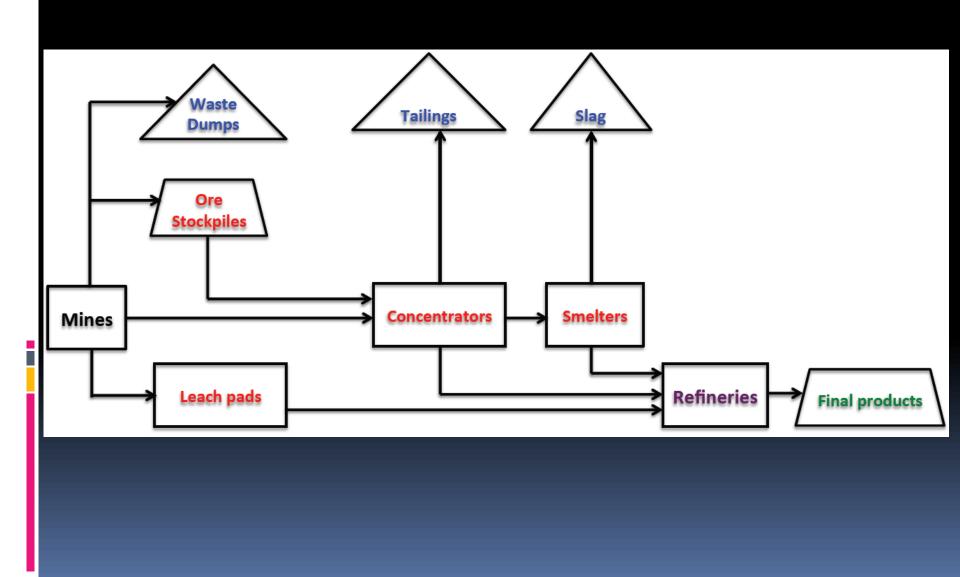
Generate set of orebody simulations

- Generate initial mine plan

<u>Loop</u>

- Modify mine plan according to metaheuristic
- Update data structures accordingly

<u>Postprocessing</u>



Closing remarks

- Mine-to-plant production scheduling is an "easy" problem, except for the supply source
- This should be reflected in the computational approach
- Resist the urge to over-centralize decisionmaking, given the geological uncertainty



Automatic Scheduling of Altonorte Operations Using Greedy Algorithms

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Oscar Mendoza Altonorte Smelter, Glencore-Xstrata



Agenda

- Introduction
- Altonorte Operations
- Greedy Programming
- Extensions/Future Work



 Common for smelters to use <u>manual</u> <u>daily scheduling</u> techniques

Two problems:

- 1. Suboptimality (w.r.t. any particular objective)
- **2.** Limits accuracy of simulation \rightarrow decicion-making

Altonorte has taken a first step

- On par with Chuquicamata (Pradenas et al., 2006)
- (we have developed more sophisticated algorithms, but not yet implemented)





- Level of Noranda Reactor : continuous
 Semi-Discrete
- Discrete PSC batches

- The Peirce-Smith Converter Problem
 - Scientific Basis for OR of Cu and Ni smelting



Dynamics

- PSC at Altonorte
 - Seven step cycle
 - **1. Initial charging** (4 to 7 ladles, depending on which converter)
 - 2. Blow
 - 3. Charge an additional ladle
 - 4. Blow
 - 5. Charge an additional ladle
 - 6. Final blow
 - 7. Final skim and discharge
 - High matte grade (~73%), so Cu-Blow dominates
 - Different sizes of converters

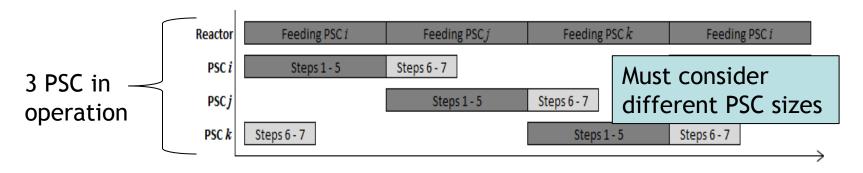


- PSC at Altonorte
 - Coordination with Noranda Reactor

- Merge steps

1-5: Initial charge, blow, charge additional ladle, blow, charge additional ladle

6-7: Final blow, skim, discharge



After Step 5, reactor is free to charge another converter

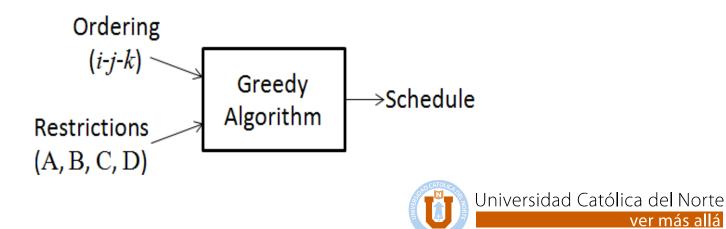


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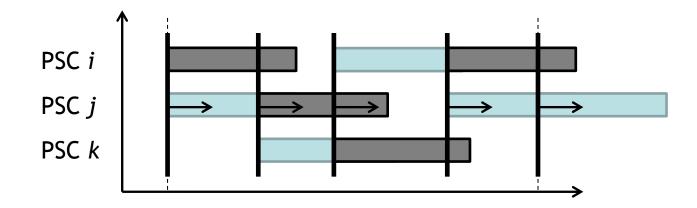
- Integrated into combinatorial search
 - Find converter sequence that maximizes number of converted ladles in schedule

Consider the following restrictions:

- A) Initial conditions
- **B)** Availability of converter
- **C)** Availability of reactor (consider production rate of reactor)
- **D) Offgas handling system**



ver más alla



- Given an ordering (*i-j-k*),
 - cursor advances from left to right
 - places the next cycle as <u>early as possible</u> (while respecting restrictions)

"Greedy" \rightarrow cursor only looks forward



- Mathematical formalism:
 - Consider ordering (*i-j-k*)
 - Consider the four restrictions A, B, C, D

$$t_1 = \text{starting time of first cycle of the new schedule} = \max\{ t_1^A, t_1^B, t_1^C, t_1^D \}$$

where t_1^{A} = earliest starting time for first cycle not to violate A t_1^{B} = earliest starting time for first cycle not to violate B t_1^{C} = earliest starting time for first cycle not to violate C t_1^{D} = earliest starting time for first cycle not to violate D

Result: *t*₁**is the earliest time which satisfies all four conditions**



- Mathematical formalism:
 - Similarly,

 $t_2 = \text{starting time of } \underline{\text{second cycle of the new schedule}} \\ = \max\{ t_2^A, t_2^B, t_2^C, t_2^D \}$

- Calculation of $(t_2^{B}, t_2^{C}, t_2^{D})$ takes into account the first cycle (thus second cycle does not conflict with <u>first cycle</u>)
- More generally,

 $t_{l} = \text{starting time of } l^{\text{th}} \text{ cycle_of the new schedule}$ $= \max\{ t_{l}^{\text{A}}, t_{l}^{\text{B}}, t_{l}^{\text{C}}, t_{l}^{\text{D}} \}$

- Calculation of $(t_l^{\rm B}, t_l^{\rm C}, t_l^{\rm D})$ takes into account the previous cycles, 1,2,3,... (*l*-1)

(thus *lth* cycle does not conflict with <u>previous cycles</u>)



Extensions/Future Work

- Paper also describes the management of <u>Refining/Casting</u>
- More advanced algorithms
 - Alternate objectives (modes of operation)
 - Convert as much matte as possible
 - Reduce a certain class of WIP
 - Use as little energy as possible, etc.

Operations Research of Cu Smelting

