

A Heuristic Approach for Short-Term Production Planning in Open Pits

Michel GAMACHE
Gabriel L'HEUREUX
François SOUMIS



GERAD

Groupe d'études et de recherche
en analyse des décisions

COSMO - Stochastic Mine Planning Laboratory
McGill University
April 18th, 2012

Definition of the problem

Short Term Planning (Tactical planning)

The short-term planning problem in open pit mines consists in scheduling mining activities over a period of 10 to 90 days while taking into account operational constraints (precedence, blending, stripping ratio, equipment capacity, etc).

For each period of the planning horizon, the schedule must indicate

- which block will be drilled, blasted and extracted;
- the location of each shovel (faces).

Objectives

- Develop an optimal approach:
 - What is the maximum size of problems that can be solved to optimality in a reasonable time?
 - What are the elements of the problem that make this problem difficult to solve?
- Identify potential heuristic approaches.

Outline

- 1 Overview of the MIP model
- 2 Strategies
- 3 Heuristic approaches

Outline

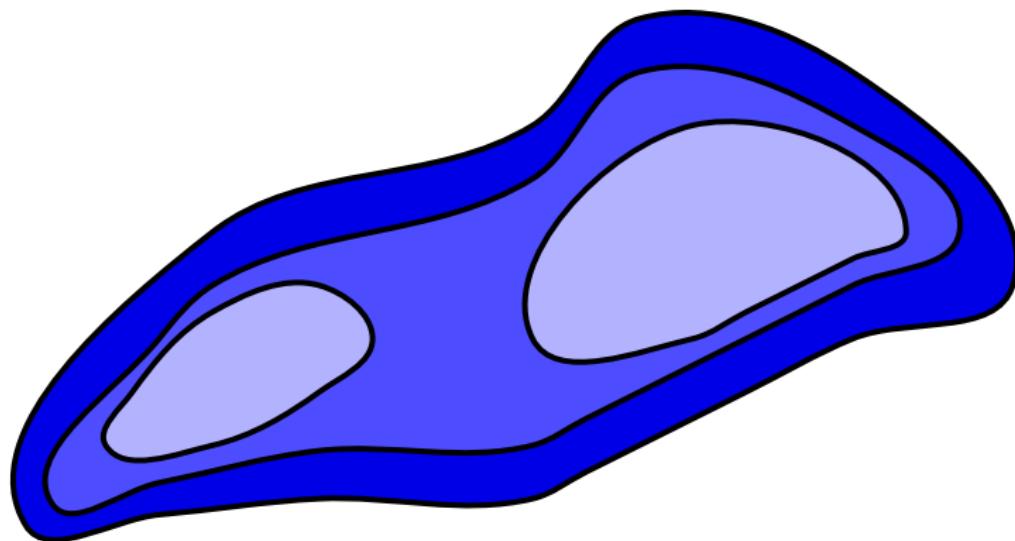
1 Overview of the MIP model

- Definitions
- Objective Function
- Constraints

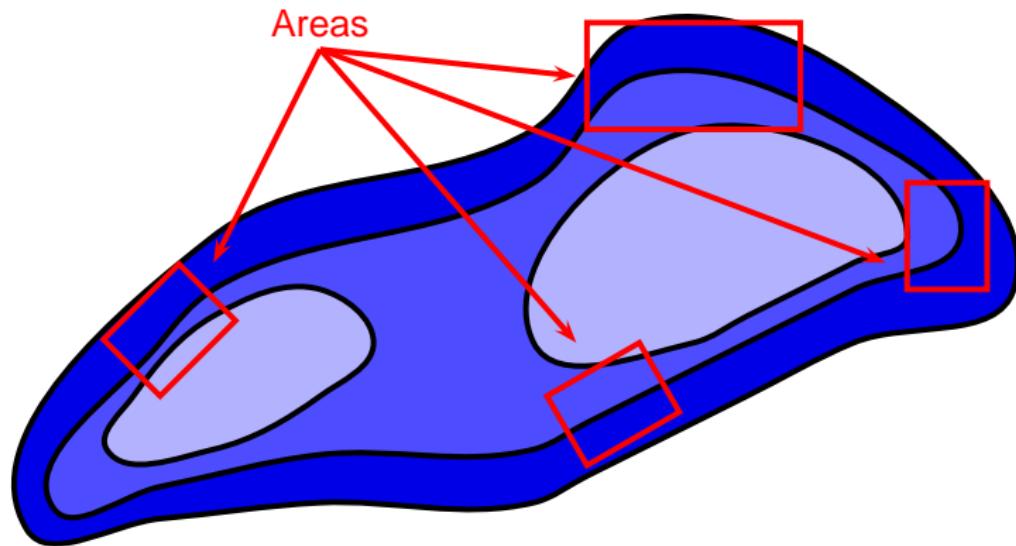
2 Strategies

3 Heuristic approaches

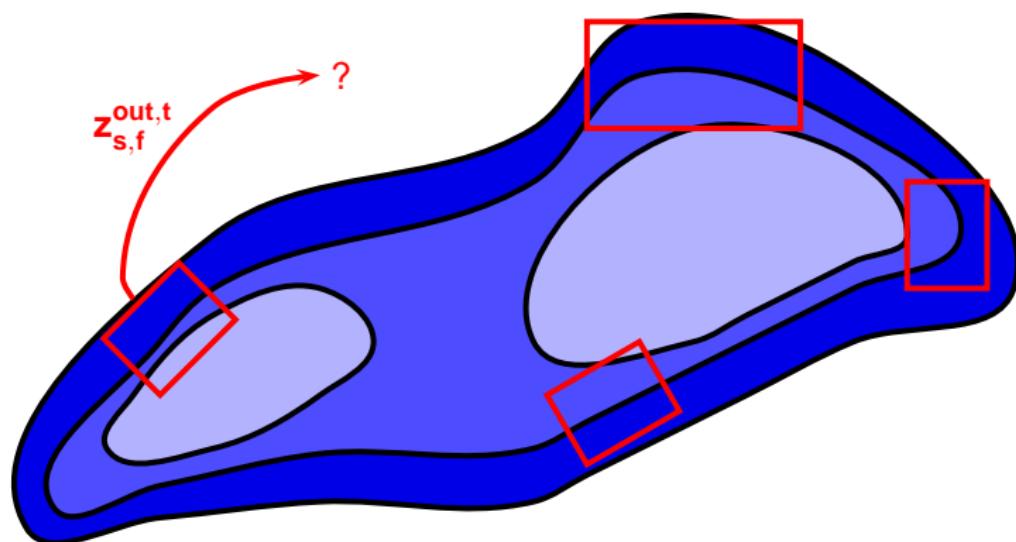
Definitions : Areas



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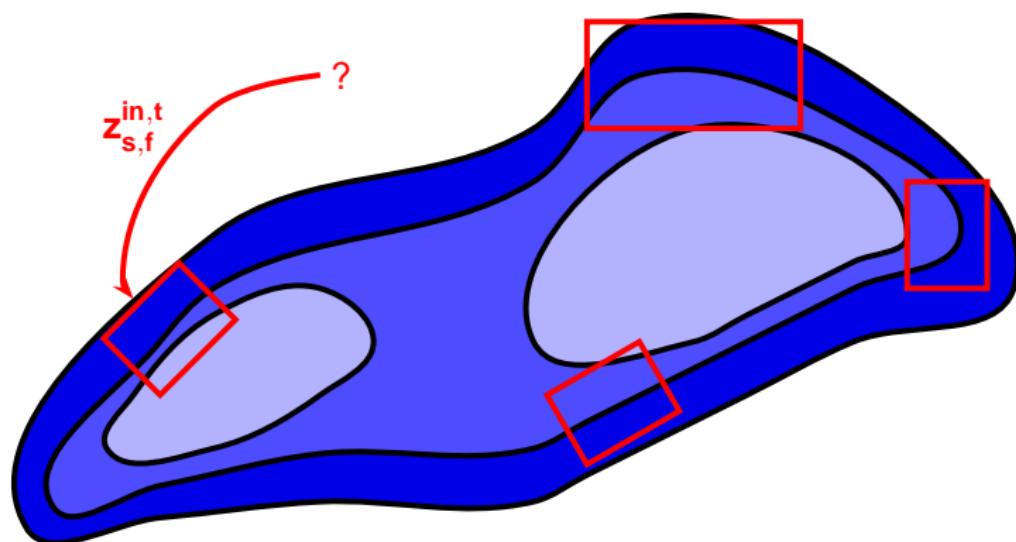


Definitions : Areas



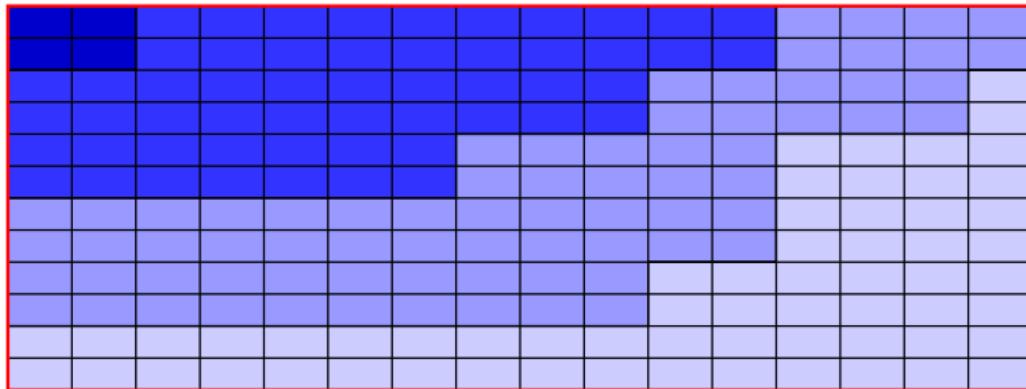
$$z_{s,f}^{out,t} = \begin{cases} 1 & \text{if shovel } s \text{ moves from face } f \text{ to another area at time } t, \\ 0 & \text{otherwise.} \end{cases}$$

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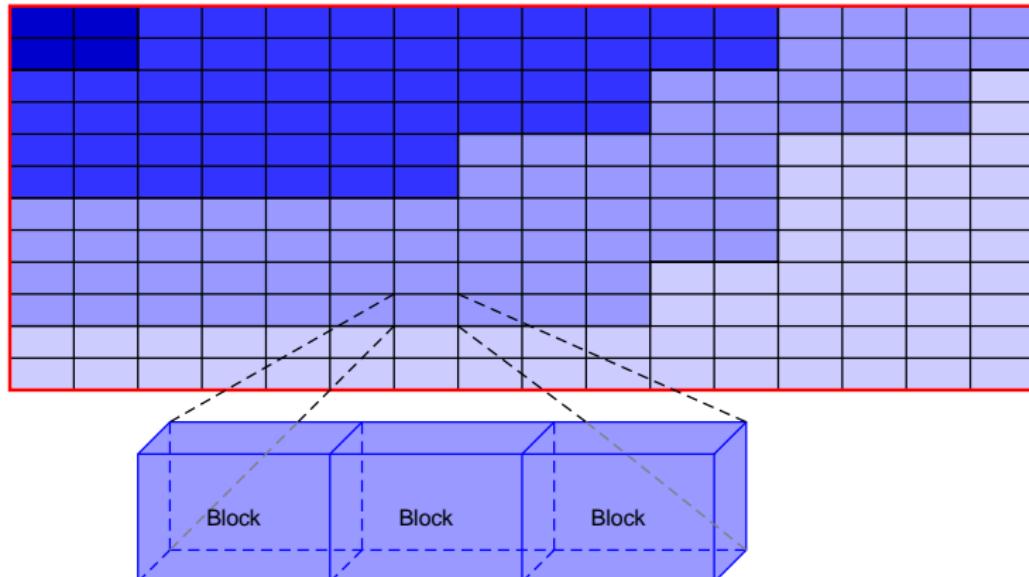


$$z_{s,f}^{in,t} = \begin{cases} 1 & \text{if shovel } s \text{ arrives at face } f \text{ from another area at time } t, \\ 0 & \text{otherwise.} \end{cases}$$

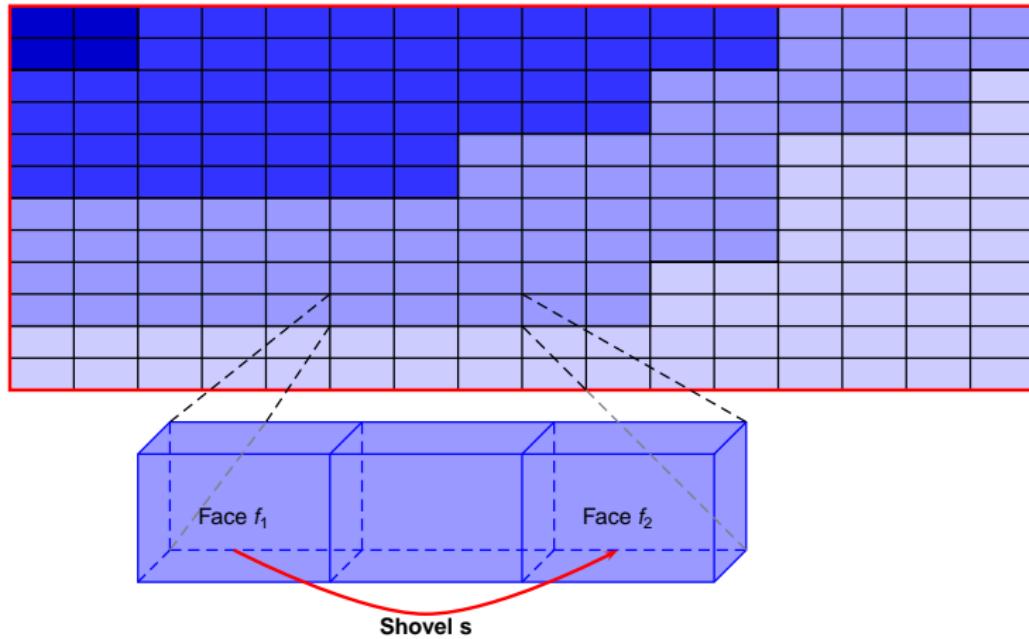
Definitions : Blocks and Faces



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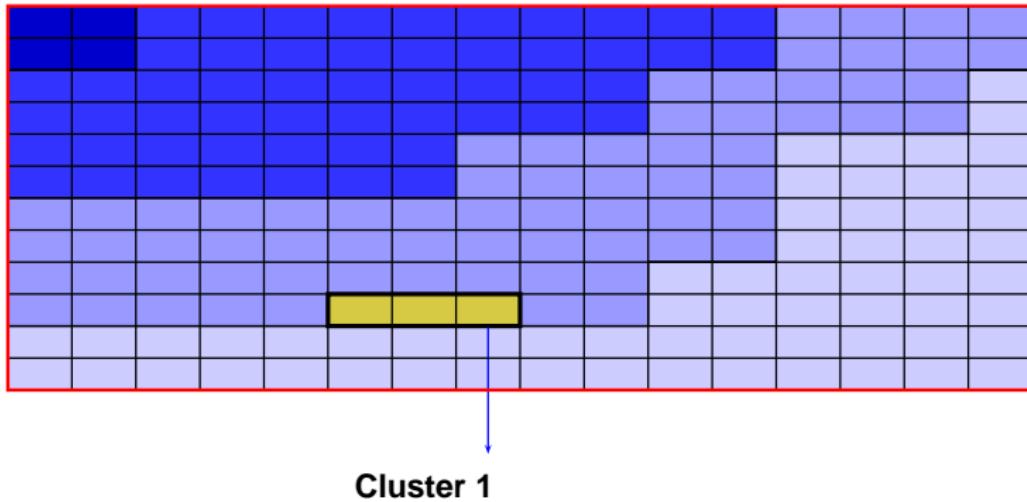


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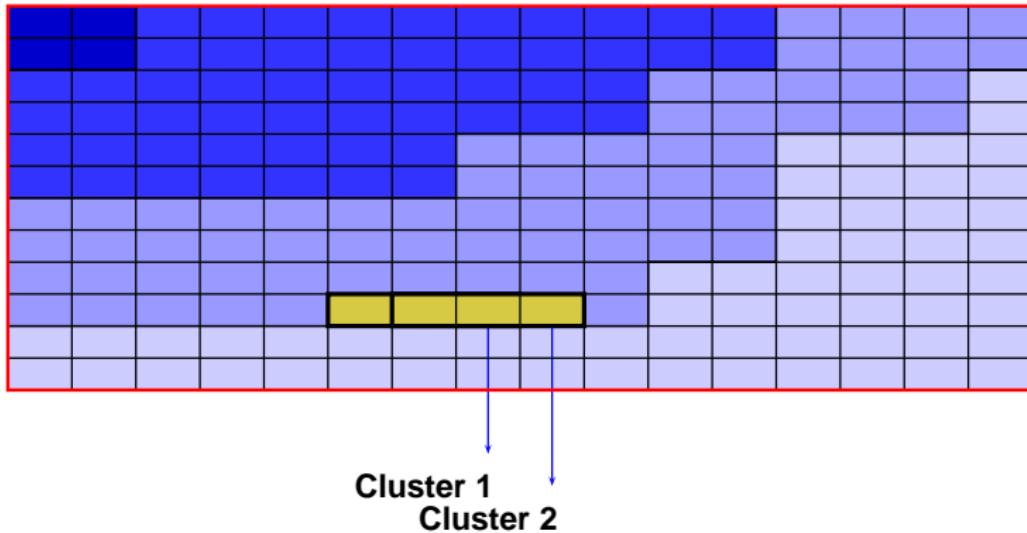


$$z_{s,f_1,f_2}^t = \begin{cases} 1 & \text{If shovel } s \text{ arrives at face } f_2 \text{ from face } f_1 \text{ at time } t, \\ 0 & \text{otherwise.} \end{cases}$$

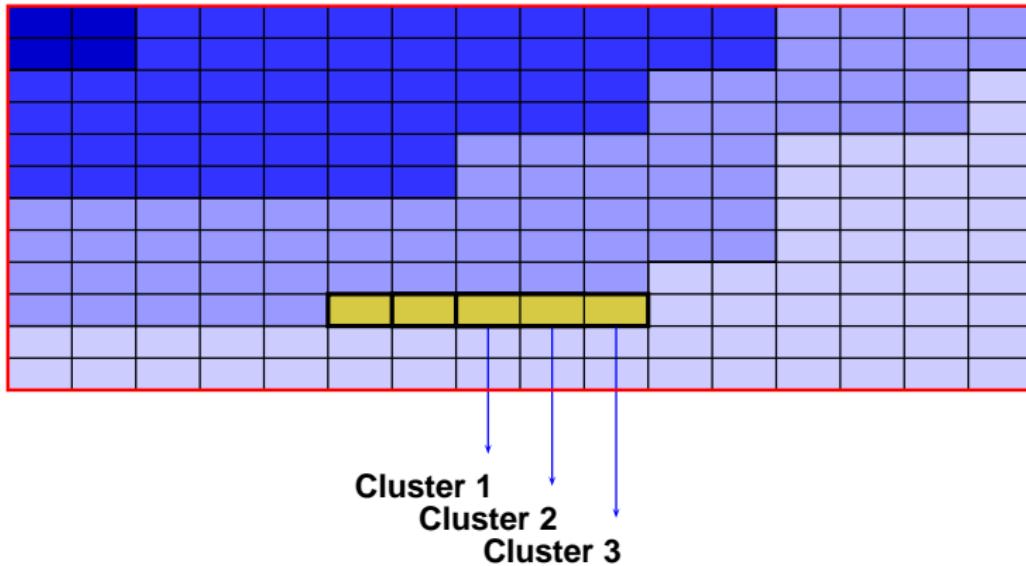
Clusters



Clusters

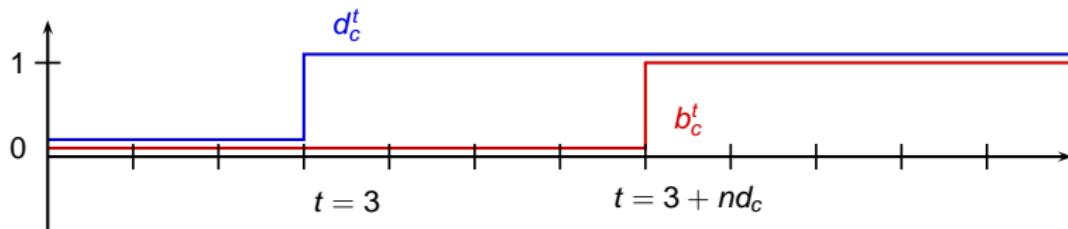


Clusters



Clusters

$$d_c^t = \begin{cases} 1 & \text{if cluster } c \text{ has been drilled at time } t, \\ 0 & \text{otherwise.} \end{cases}$$
$$b_c^t = \begin{cases} 1 & \text{if cluster } c \text{ has been blasted at time } t, \\ 0 & \text{otherwise.} \end{cases}$$



Objective function

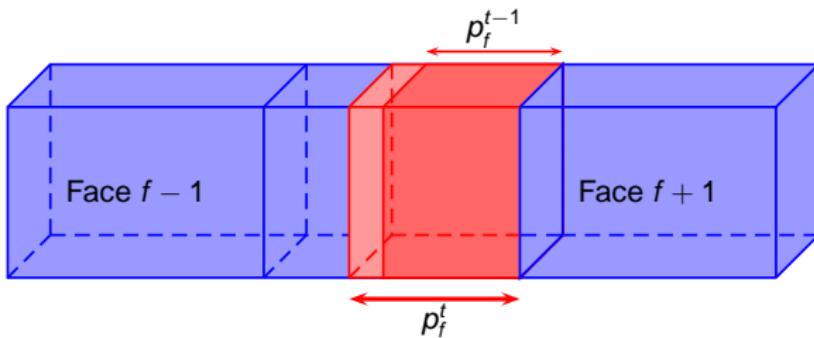
$$\begin{aligned}
 & \min \sum_s \sum_t \sum_{\forall f_1 \in F} TCU(z_{s,f_1}^{in,t} + z_{s,f_1}^{out,t}) && (\textit{Long-distance movements}) \\
 & + \sum_s \sum_t \sum_{\forall f_1 \in F} \sum_{\forall f_2 \in F | f_2 \neq f_1} \sum_{T} TC_{f_1, f_2} * z_{s,f_1,f_2}^t && (\textit{Movements within the same area}) \\
 & + \sum_c (DC_c^1 d_c^1 + \sum_{t=2}^T DC_c^t (d_c^t - d_c^{t-1})) && (\textit{Drilling}) \\
 & + \sum_c (BC_c^1 b_c^1 + \sum_{t=2}^T BC_c^t (b_c^t - b_c^{t-1})) && (\textit{Blasting})
 \end{aligned}$$

Constraints

Blending and Stripping Ratio

$$\sum_{f=1}^F m_f \alpha_f q_f^e (p_f^t - p_f^{t-1}) \geq L_e^t \sum_{f=1}^F m_f \alpha_f (p_f^t - p_f^{t-1}) \quad t = 2, \dots, T \quad \forall e \quad (2)$$

$$\sum_{f=1}^F m_f \alpha_f q_f^e (p_f^t - p_f^{t-1}) \leq U_e^t \sum_{f=1}^F m_f \alpha_f (p_f^t - p_f^{t-1}) \quad t = 2, \dots, T \quad \forall e \quad (3)$$



Constraints

Shovels' location and production capacity

$$w_{s,f}^t = \begin{cases} 1 & \text{if shovel } s \text{ is assigned to face } f \text{ during period } t, \\ 0 & \text{otherwise.} \end{cases}$$

$$\sum_f w_{s,f}^t \leq 1 \quad \forall t; \forall s \quad (4)$$

$$\sum_s w_{s,f}^t \leq 1 \quad \forall t; \forall f \quad (5)$$

$$\sum_f m_f (p_f^t - p_f^{t-1}) \geq LS \sum_{s=1}^S Us_s \quad t = 2, \dots, T \quad (6)$$

$$m_f (p_f^t - p_f^{t-1}) \leq \sum_s Us_s w_{s,f}^t \quad t = 2, \dots, T, \forall f \quad (7)$$

Constraints

Resource usage at time t

Number of drilling

$$\sum_c (d_c^t - d_c^{t-nd_c}) \geq Ld^t \quad t = 2, \dots, T \quad (8)$$

$$\sum_c (d_c^t - d_c^{t-nd_c}) \leq Ud^t \quad t = 2, \dots, T \quad (9)$$

Number of blasting

$$\sum_c (b_c^t - b_c^{t-1}) \geq Ld^t \quad t = 2, \dots, T \quad (10)$$

$$\sum_c (b_c^t - b_c^{t-1}) \leq Ud^t \quad t = 2, \dots, T \quad (11)$$

Constraints

The number of faces in operation at time t

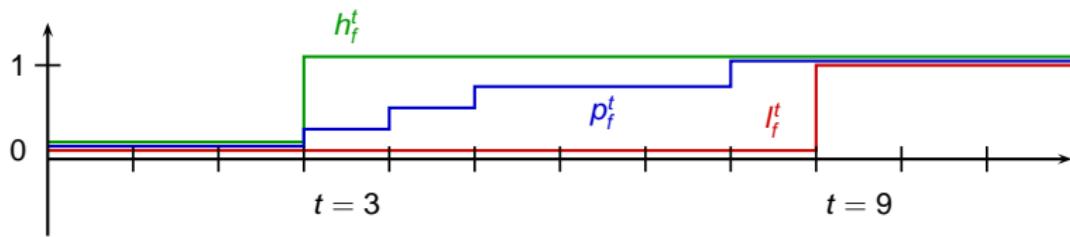
$$h_f^t = \begin{cases} 1 & \text{if the extraction of the material at face } f \text{ has begun at time } t, \\ 0 & \text{otherwise.} \end{cases}$$

$$l_f^t = \begin{cases} 1 & \text{if all the material at face } f \text{ has been extracted at time } t, \\ 0 & \text{otherwise.} \end{cases}$$

$$\sum_f (h_f^t - l_f^t) \leq nf \quad t = 2, \dots, T \quad (12)$$

$$l_f^t \leq p_f^t \quad \forall t; \forall f \quad (13)$$

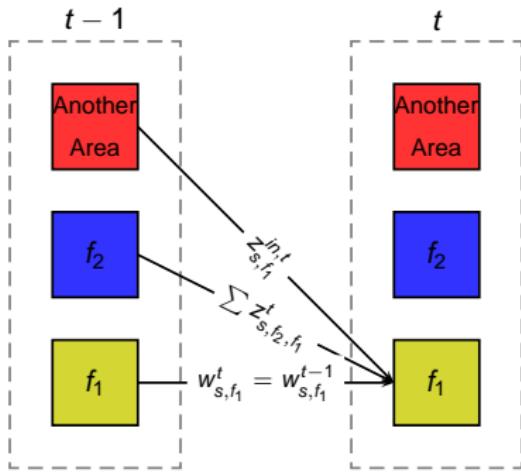
$$p_f^t \leq h_f^t \quad \forall t; \forall f \quad (14)$$



Constraints

Flow conservation

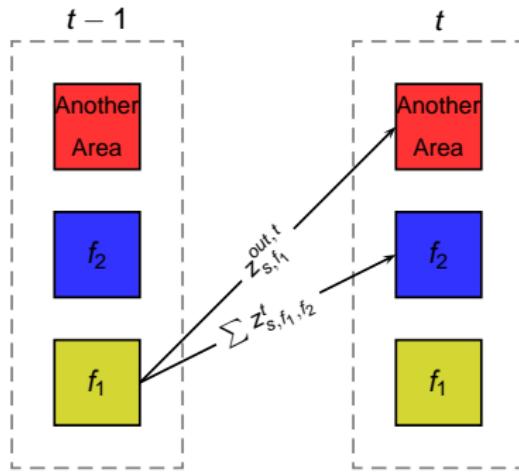
$$w_{s,f_1}^t = w_{s,f_1}^{t-1} + \left(z_{s,f_1}^{in,t} + \sum_{f_2 \in ARCS_{f_1}} z_{s,f_2,f_1}^t \right) - \left(z_{s,f_1}^{out,t} + \sum_{f_2 \in ARCS_{f_1}} z_{s,f_1,f_2}^t \right) \quad t = 2, \dots, T; \forall s \forall f_1 \in F \quad (15)$$



Constraints

Flow conservation

$$\begin{aligned} w_{s,f_1}^t = w_{s,f_1}^{t-1} &+ \left(z_{s,f_1}^{in,t} + \sum_{f_2 \in ARCS_{f_1}} z_{s,f_2,f_1}^t \right) \\ &- \left(z_{s,f_1}^{out,t} + \sum_{f_2 \in ARCS_{f_1}} z_{s,f_1,f_2}^t \right) \quad t = 2, \dots, T; \forall s \forall f_1 \in F \end{aligned} \quad (15)$$



Constraints

Precedence constraints

- Precedence between operations

Drilling → Blasting → Excavation

$$d_c^{t-n_{dc}} \geq b_c^t \quad t = 2, \dots, T; \forall c \quad (16)$$

$$\sum_{c \in C_f} b_c^{t-n_{bc}} \geq h_f^t \quad t = 2, \dots, T; \forall f \quad (17)$$

$$l_f^{t-1} \geq b_c^t \quad t = 2, \dots, T; \forall c; \forall f \in F_c \quad (18)$$

- Temporal link

$$d_c^{t-1} \leq d_c^t \quad t = 2, \dots, T; \forall c \quad (19)$$

$$b_c^{t-1} \leq b_c^t \quad t = 2, \dots, T; \forall c \quad (20)$$

$$h_f^{t-1} \leq h_f^t \quad t = 2, \dots, T; \forall f \quad (21)$$

$$l_f^{t-1} \leq l_f^t \quad t = 2, \dots, T; \forall f \quad (22)$$

$$p_f^{t-1} \leq p_f^t \quad t = 2, \dots, T; \forall f \quad (23)$$

Constraints

Clusters

$$y_c = \begin{cases} 1 & \text{if cluster } c \text{ is chosen,} \\ 0 & \text{otherwise.} \end{cases}$$

$$\sum_{c \in K_k} y_c \leq 1 \quad \forall k \quad (24)$$

$$d_c^T \leq y_c \quad \forall c; \quad (25)$$



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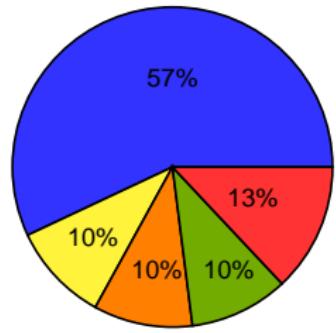
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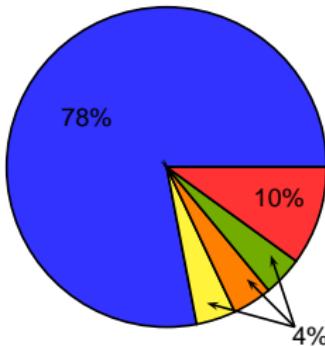
Instances

	T1	T2	T3
Periods	10	30	90
Faces	30	90	192
Areas	5	5	8
Shovels	5	5	5
Clusters	60	315	864
Variables	15 660	310 815	2 541 024
Constraints	5 912	83 729	627 178

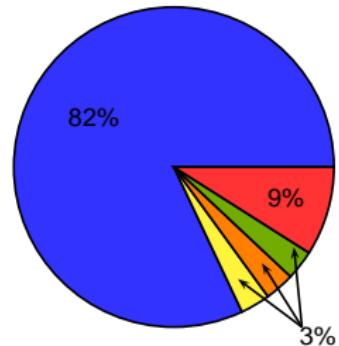
Distribution of variables in the models



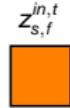
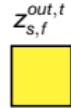
T1



T2



T3



Results : M_0 (basic model)

	Number Constraints	Number Variables	Number Nodes	Integrality Gap (%)	Times Z_{IP}^{opt} (sec)	Times $Z_{IP}^{5\%}$ (sec)	Times Z_{IP}^{Best} (sec)
M_0	4060	12020	1457910	N/A	36000	N/A	N/A

Results : M_1

- ➊ Set the value of some variables based on the earliest start of each task (considering precedence constraints only);
- ➋ Prioritize the branching on the variables $w_{s,f}^t$.

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M_1	2680	10735	736204	31%	6464	40	40

- ➊ At least 5 times faster;
- ➋ The number of nodes was reduced by half;
- ➌ The best solution found in less than 40 sec.

Results : M_2

- ➊ Symmetric solutions;
- ➋ Adding discounts in the objective function based on the ore grade of each face : $-\sum_t \sum_f r_f^t p_f^t$;

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M_2	2423	8875	24682	33%	200	5	17

- At least 30 (180) times faster than M1 (M_0);
- 10% reduction on the number of constraints and variables;
- 97% reduction on the branching nodes;
- Integrality gap ???

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Results : M_3

- ➊ Fractional variables ($w_{f,s}^t$) \Rightarrow this eliminates almost all the movements of shovels which greatly reduces costs;
- ➋ Solution \Rightarrow Adding a new constraint that forces a minimum number of movements for each shovel.

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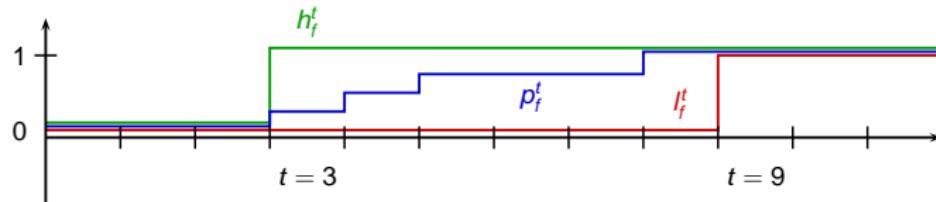
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M_2	2423	8875	24682	33%	200	5	17
M_3	2664	10557	38687	0.8%	400	13	310

- ➊ It reduces the gap;
- ➋ The solution time increases.

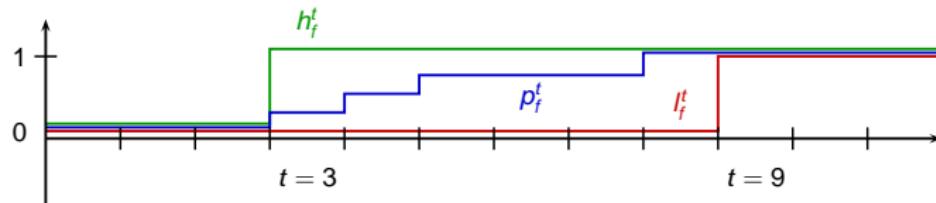
Results : M_4

Solution of the MILP :

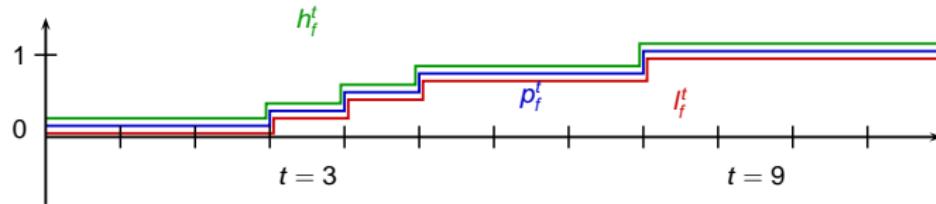


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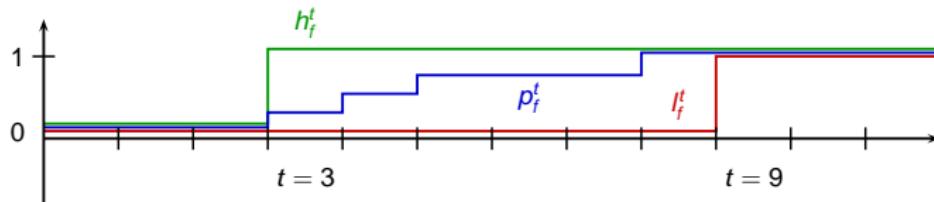


Solution of the LP relaxation

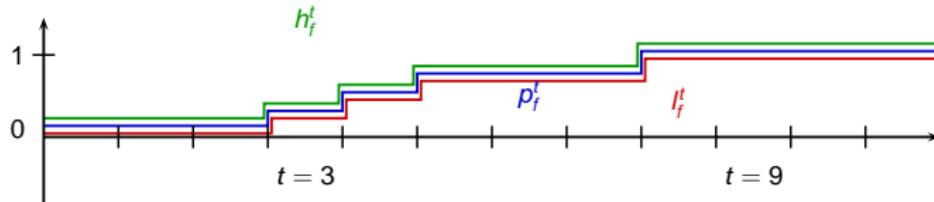


Results : M_4

Solution of the MILP :



Solution of the LP relaxation



$$1 - (h_f^t - h_f^{t-1}) \geq l_f^{t+1} \quad t = 2, \dots, T-1; \forall f \quad (41)$$

$$1 - (h_f^t - h_f^{t-1}) \geq l_f^{t+2} \quad t = 2, \dots, T-2; \forall f \quad (42)$$

Results : M_4

- ➊ Based on M2
- ➋ Stronger link between variables :

$$w_{s,f_1}^t \geq \sum_{\forall f_2 \in ARCS_{f_1}} z_{s,f_2,f_1}^t \quad \forall t; \quad \forall f_1 \in F; \quad \forall s \quad (39)$$

$$w_{s,f_1}^t \geq \sum_{\forall f_3 \in ARCS_{f_1}} z_{s,f_1,f_3}^t \quad t = 1, \dots, T-1; \quad \forall f_1 \in F; \quad \forall s \quad (40)$$

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$$1 - (h_f^t - h_f^{t-1}) \geq l_f^{t+2} \quad t = 2, \dots, T-2; \quad \forall f \quad (42)$$

$$d_c^t = b_c^{t+3} \quad t = 1, \dots, T-3; \quad \forall c \quad (43)$$

$$h_f^t = l_f^{t+5} \quad t = 1, \dots, T-5; \quad \forall c \quad (44)$$

$$y_c = d_c^T \quad \forall c \quad (45)$$

$$y_c = b_c^T \quad \forall c \quad (46)$$

Results : M_4

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M_4	1718	830	22145	33%	31	0	9

- 6 (1500) times faster than M2 (M0);
- 90% less variables than M2;
- 50% less constraints than M0;
- 98% less nodes in the B&B than M0.

Outline

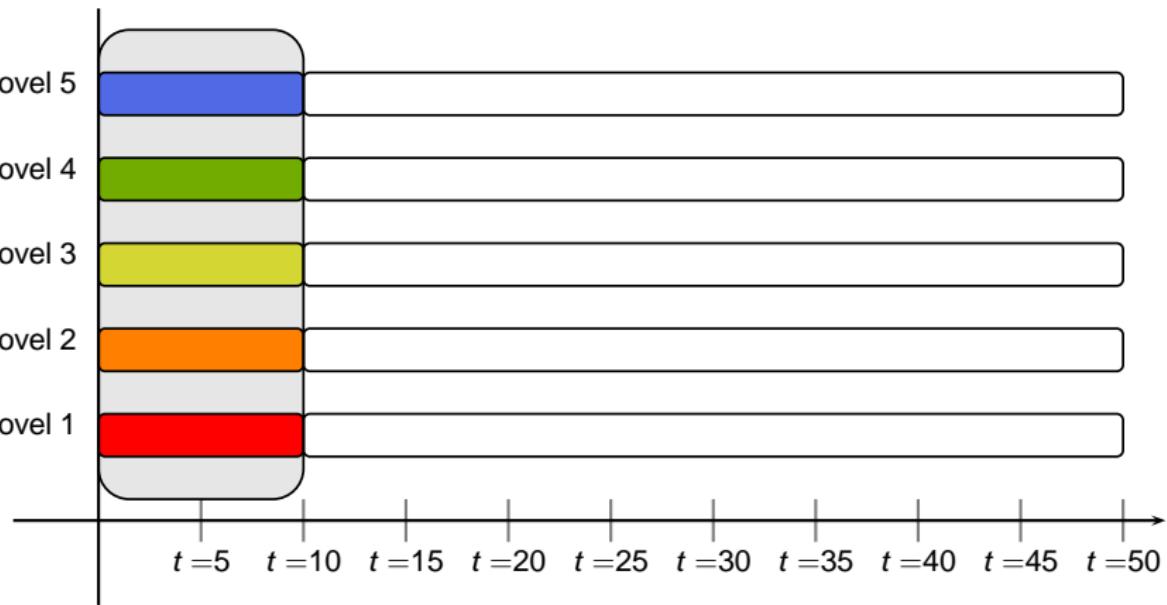
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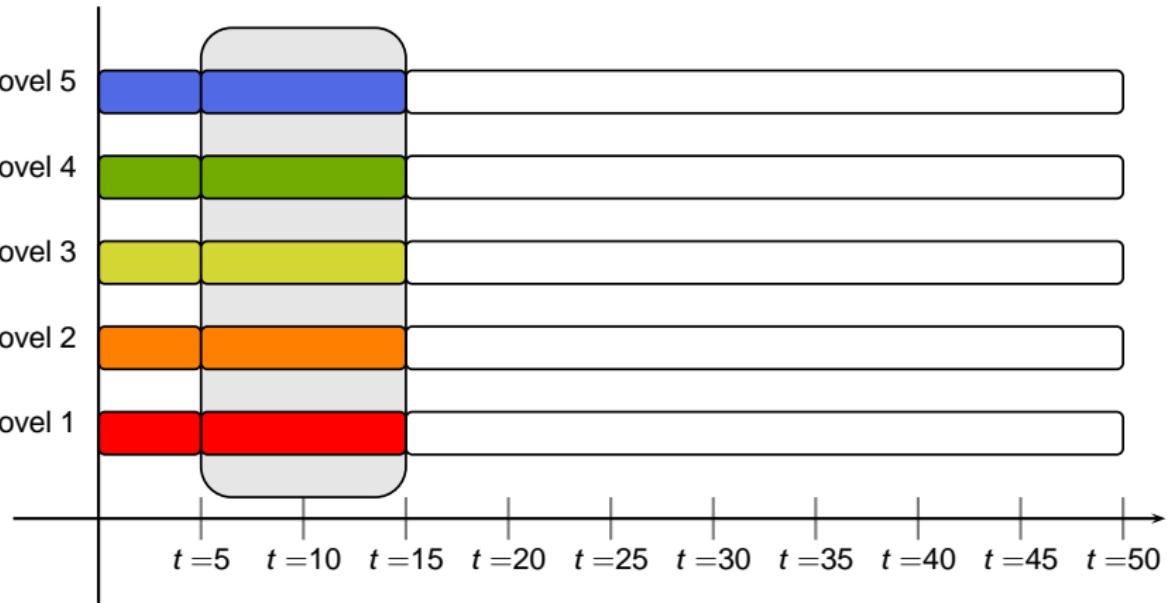
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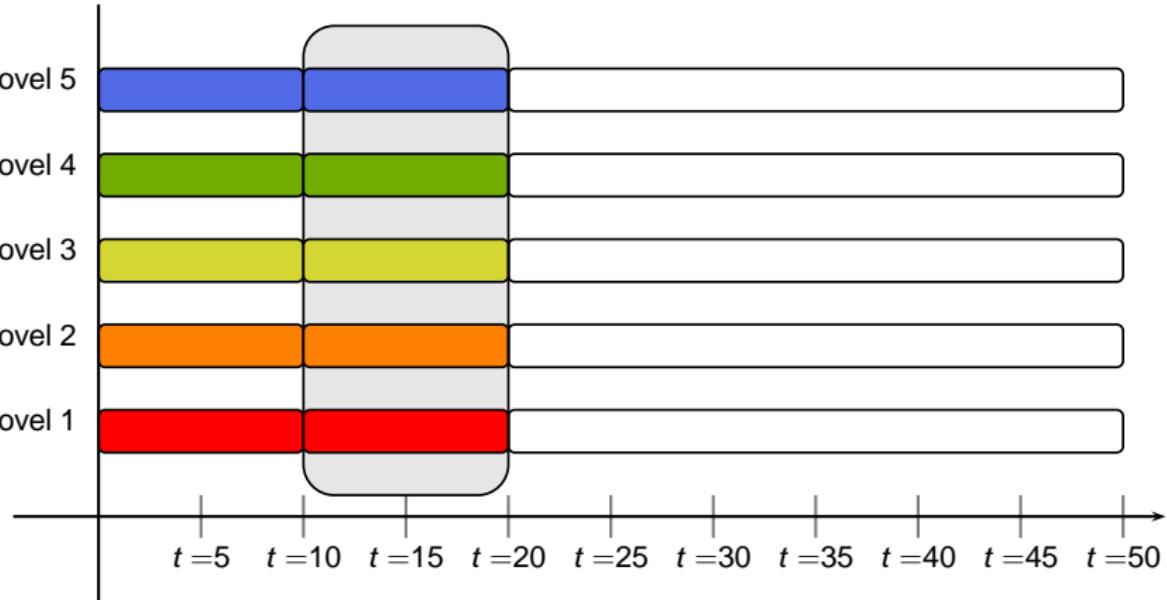
Sliding Time Horizon



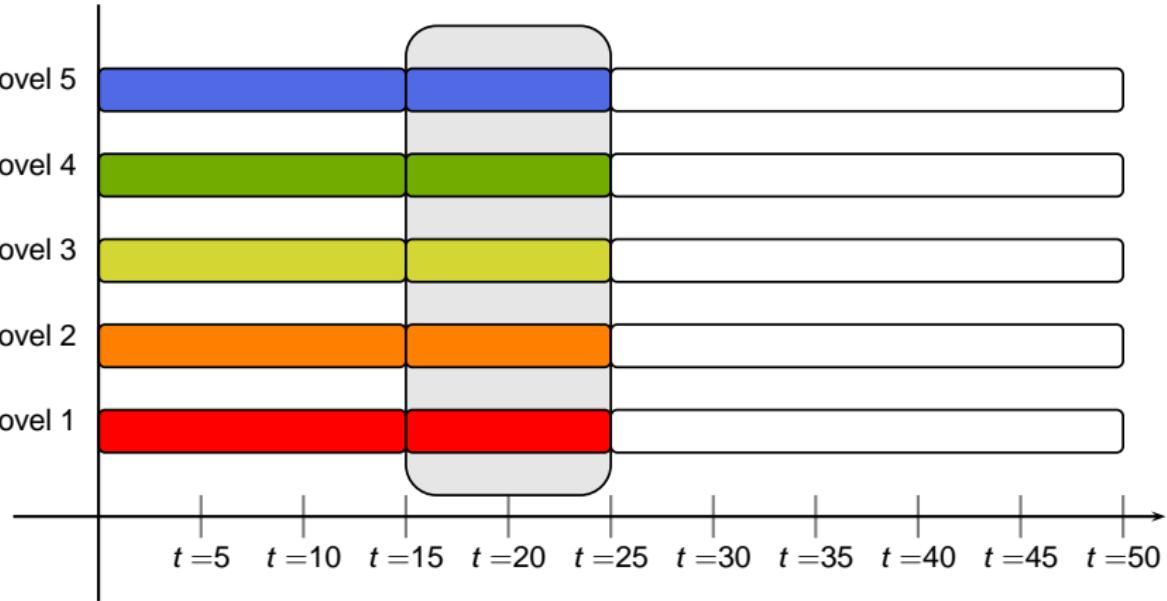
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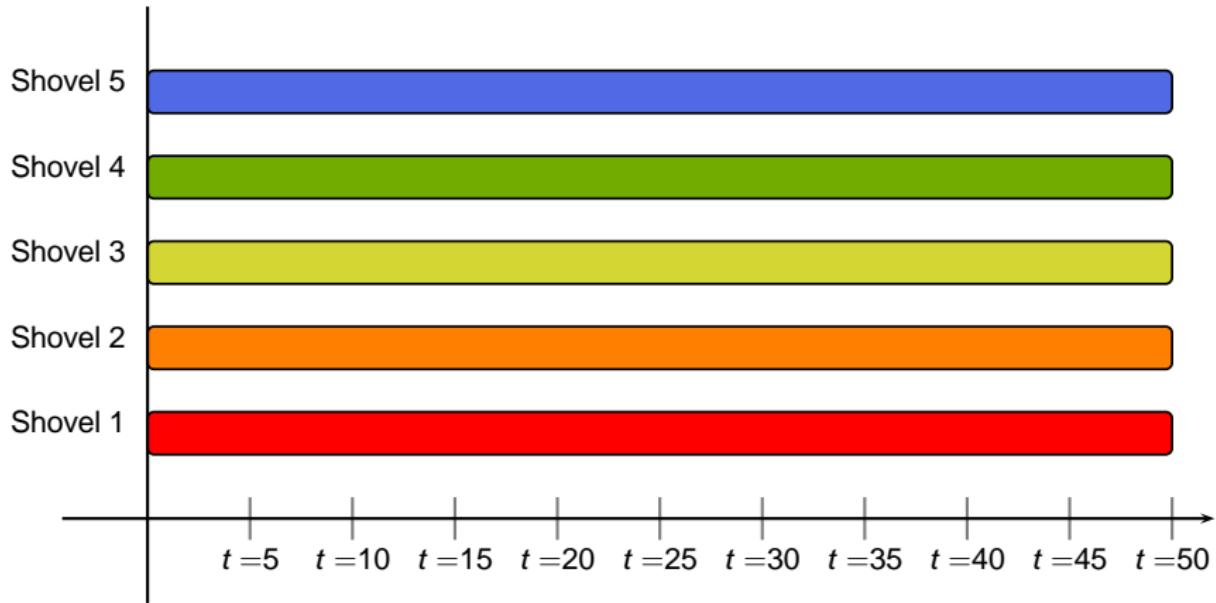
Sliding Time Horizon



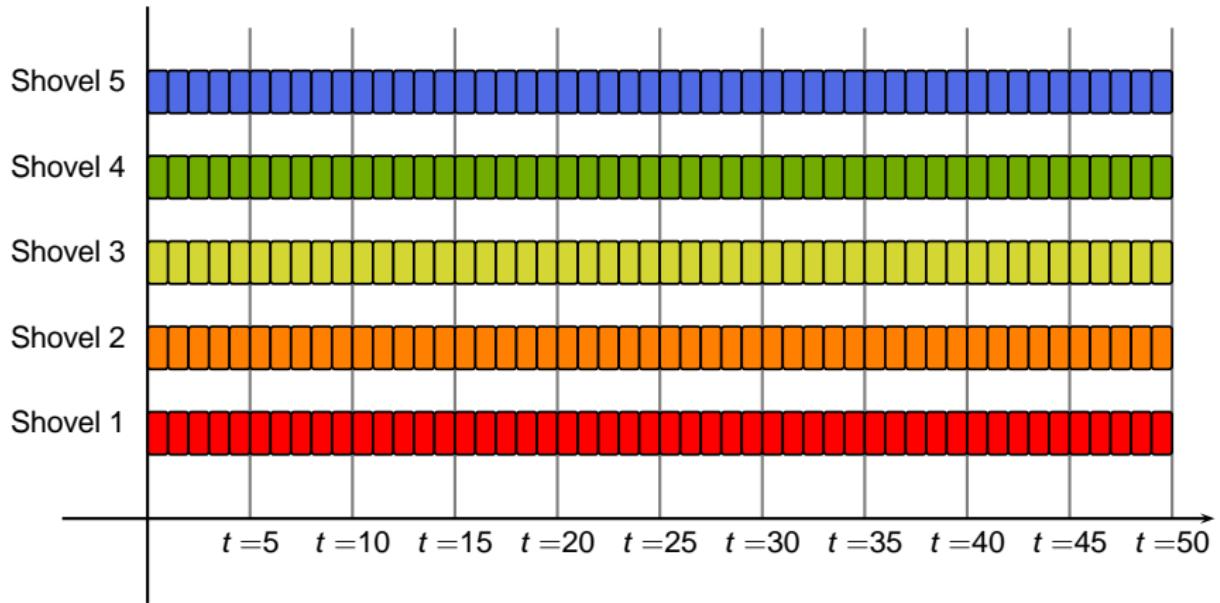
Sliding Time Horizon



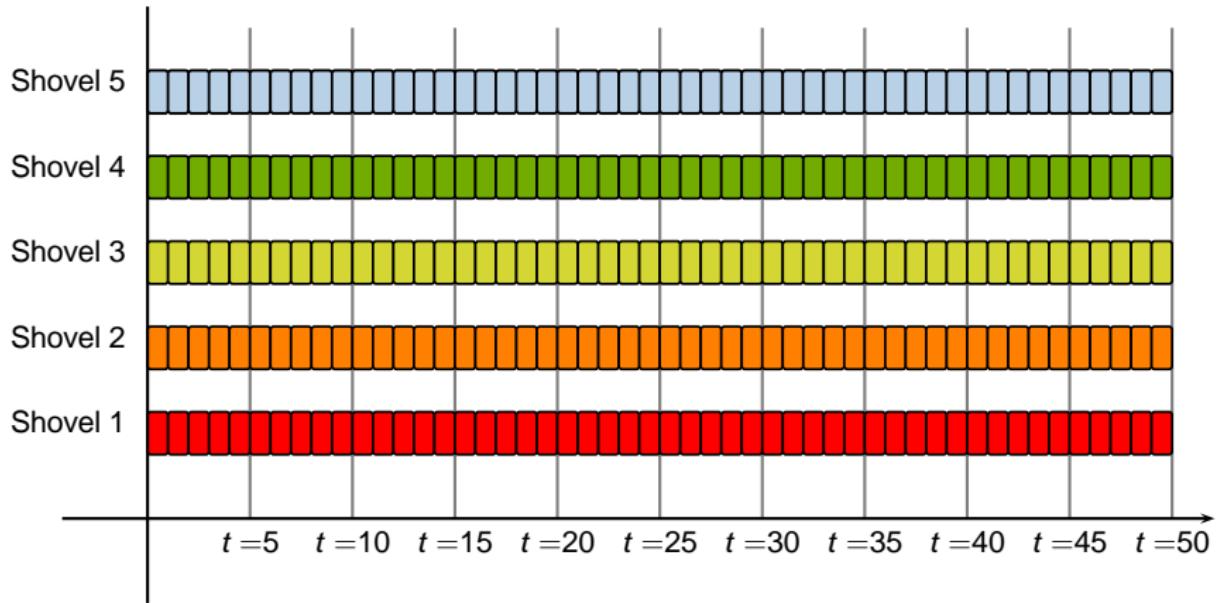
Sliding Time Horizon



Large Neighborhood Search



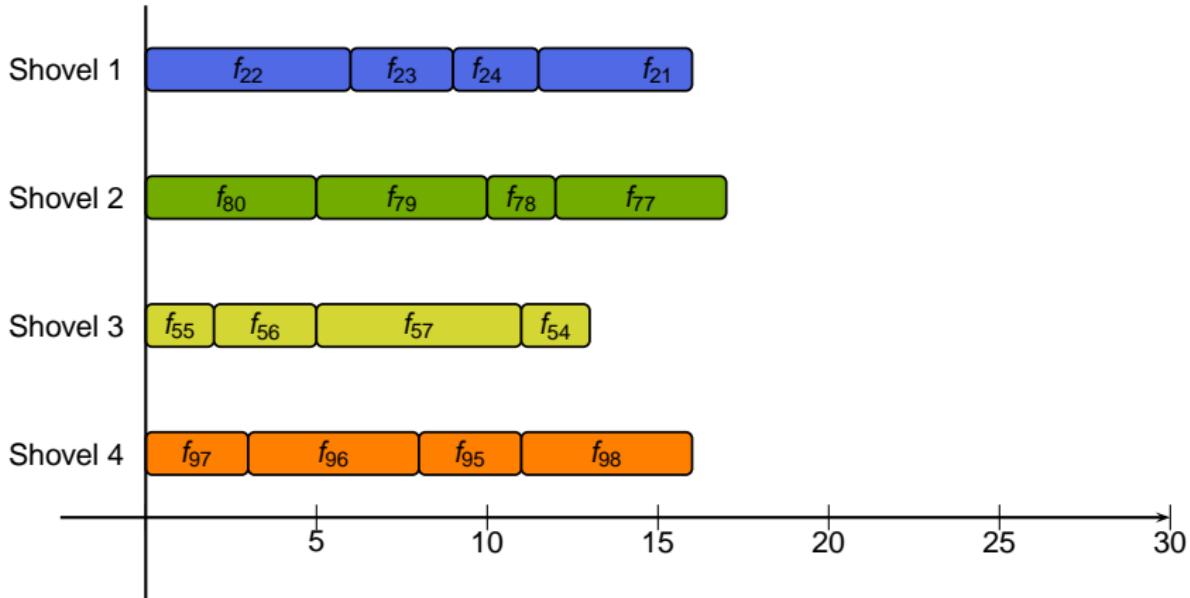
Large Neighborhood Search



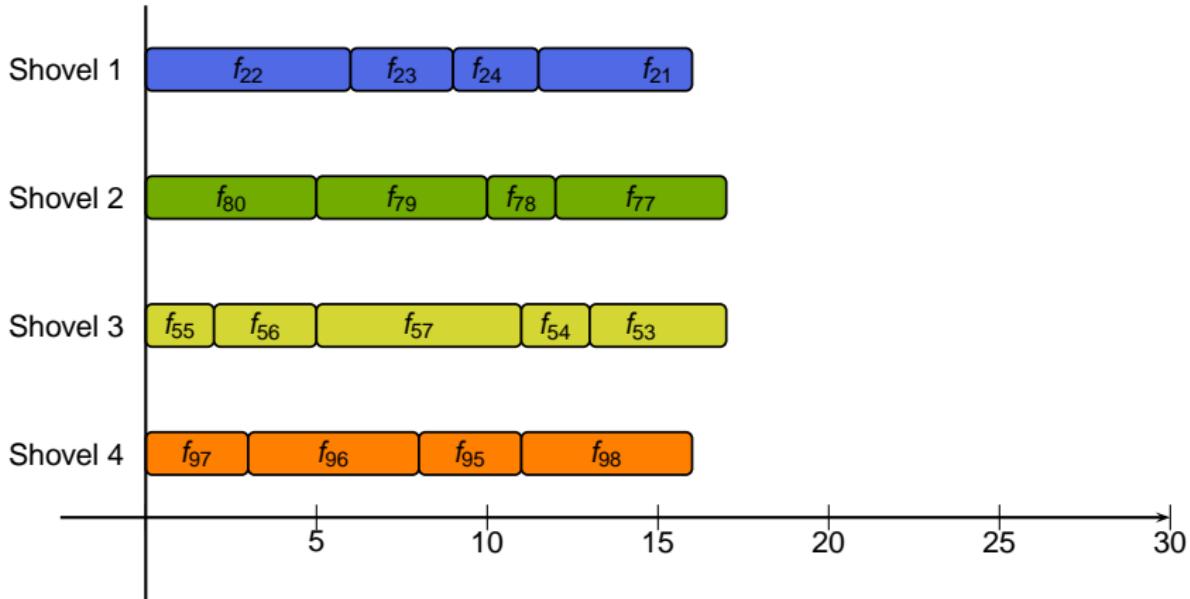
Heuristic Approach

- Construct a simple and intuitive schedule;
- Use the heuristic schedule as a guide when solving model M4

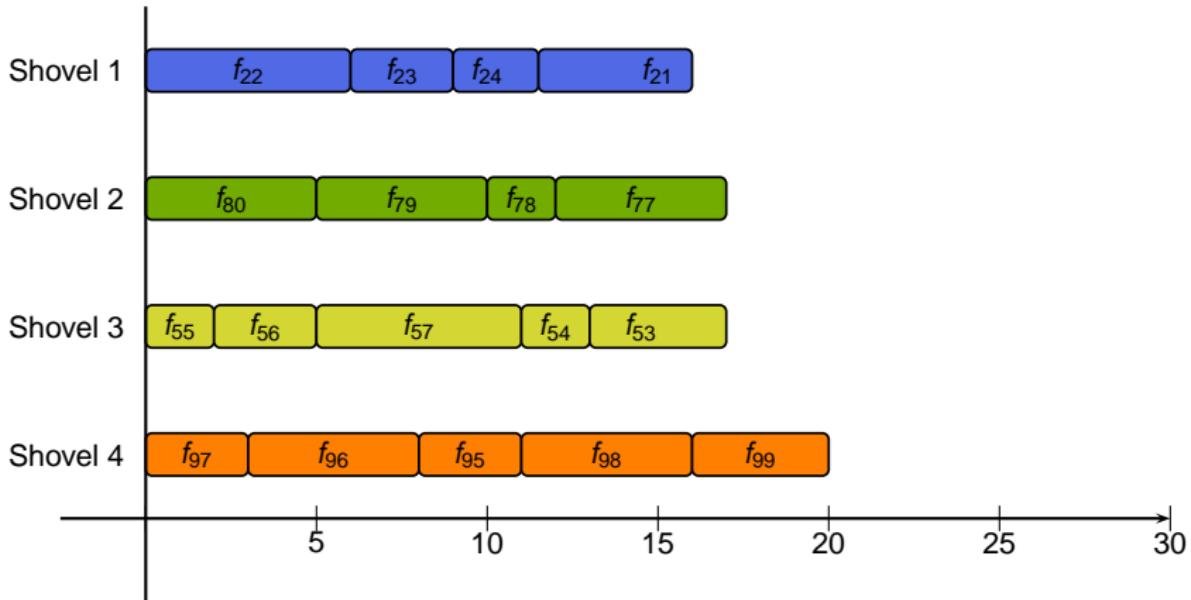
The heuristic schedule



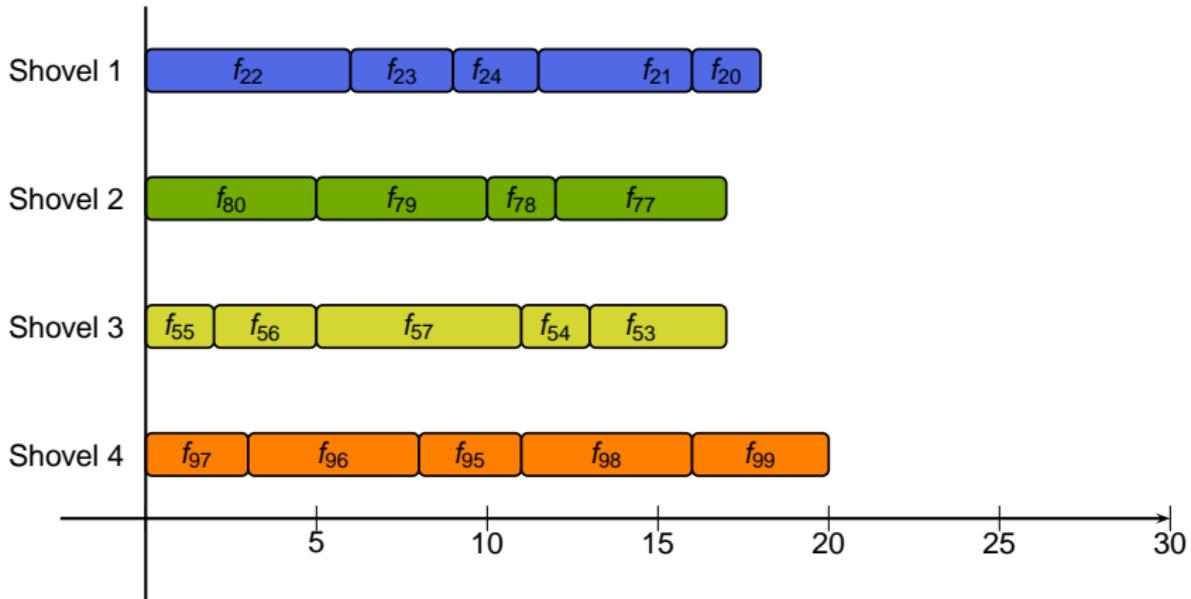
The heuristic schedule



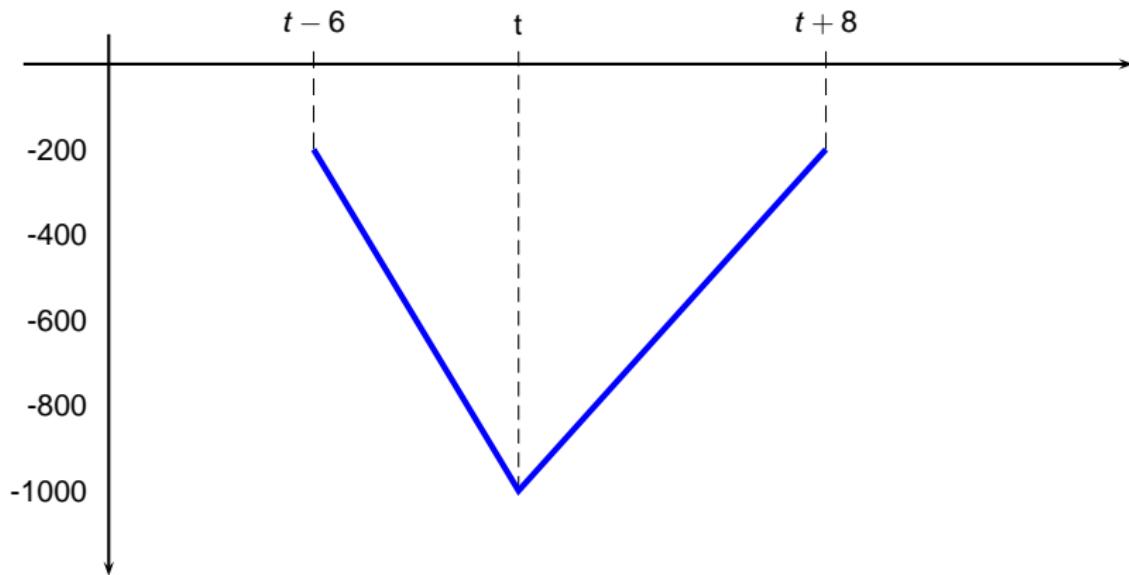
The heuristic schedule



The heuristic schedule



Reduction



Tests

- Instance T1:

	Number Constraints	Number Variables	Number Nodes	Times Z_{IP}^{opt} (sec)
M_4	1718	830	22145	31
Heuristic	1765	1290	494	9

- Instance T2:

- 30 periods;
- 90 faces;
- 310000 variables;
- 84000 constraints.

Solved in 29 seconds (4 seconds if some variables are set to zero)